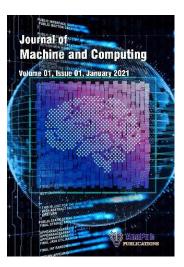
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# Advancing Security and Scalability: A Protocol Extension for Dynamic Group Membership Management

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#### Abstract

The integration of Contributory Group Key Agreement (CGKA) for group h ation revolutionizes the collaborative process of generating group keys, instilling trust and fostering collabor on along group members. By ensuring that each member actively contributes to the generation of the gree GKA distributes the responsibility of key generation across the group, thereby enhancing the security flience of the group's cryptographic and . infrastructure. Concurrently, the utilization of Lattice Diffie-H or key generation leverages the (DH) mathematical properties of lattices to securely derive shared sec ers a robust and efficient method for keys. I generating keys in cryptographic applications, ensuri tiality and integrity of communication channels. nfið Furthermore, the incorporation of blockchain penting membership changes introduces a nology br imp decentralized and transparent approach to mana membership dynamics. By leveraging blockchain's g gr ership changes can be executed securely, transparently, and distributed ledger technology and smart contracts, me efficiently. This enhances the integrity and resilience of group's membership management system, allowing for the secure addition and removal of members from the group while maintaining the integrity of the cryptographic of CCKA, LDH, and blockchain technology presents a comprehensive infrastructure. Together, the integration solution for advancing the security dynamic group membership management protocols, offering a nity robust framework for secure and icient computation in contemporary environments. Moreover, the proposed integration of CGKA, LDH, and blo chain technology facilitates seamless adaptation to dynamic changes in group membership, ensuring that and calability are maintained even as the composition of the group evolves. e evaluations, the effectiveness of the integrated approach that is implemented Through simulations and erforman in Python Software is dem strated ompared to existing protocols like Elliptic Curve Diffie-Hellman (ECDH), RSA Quantum Cryptography (PQC). Key Exchan

Keywork Conclutory Group Key Agreement, Lattice Diffie-Hellman, Blockchain, Group Membership Managemen and Security

# 1. Introduction

nami group membership management is a critical facet of contemporary communication and collaboration vstems, particularly in the realm of distributed computing, cloud computing, and decentralized networks. In such environments, groups are not static entities; instead, they are subject to frequent changes in membership due to various factors such as user additions, departures, role changes, or system failures [1] [2]. Ensuring the seamless integration and operation of new members while maintaining the security and integrity of group communication channels poses significant challenges to system designers and administrators. The traditional approach to group membership management often involves centralized systems where a single authority is responsible for managing membership changes. However, such centralized systems are inherently limited in their scalability, fault tolerance, and susceptibility to single points of failure. Moreover, they may not be well-suited for distributed or decentralized environments where autonomy, resilience, and privacy are paramount. As a result, there has been a growing interest in developing decentralized and distributed protocols for dynamic group membership management.

In dynamic group membership management, the primary objective is to facilitate the seamless addition and removal of members from a group while preserving the security, confidentiality, and integrity of group communication [3] [4]. This involves not only managing access control and authentication mechanisms but also ensuring robustness and resilience of cryptographic protocols used for key distribution, encryption, and authentic ion. Moreover, dynamic group membership management protocols must be capable of adapting to changing dynamics in real-time without compromising security or performance [14] [15]. One of the key challeng group membership management is achieving consensus among group members regarding membership le nanges v preventing unauthorized access or malicious activities [5] [6]. Traditional cryptographic technic ickey infrastructure (PKI) or shared secret key schemes may not be sufficient to address ha ges, est ally in large-scale distributed systems where the number of participants is constantly change e is a need for . The ore. innovative approaches that combine cryptographic primitives, distributed conse ims, and decentralized s algo governance mechanisms to ensure the security and scalability of dynamic group me p management protocols.

In recent years, advancements in blockchain technology, cryptographic primitives as threshold signatures and multi-party computation, and distributed consensus algorithms have paved the for new approaches to dynamic group membership management [16][17]. These approaches leverage nt properties of blockchain, such as 1111 decentralization, transparency, and immutability, to securely manage nen ership changes without relying on centralized authorities [18] [19]. Additionally, they utilize cryp techr ques to ensure the confidentiality, grap in the presence of malicious actors or network integrity, and authenticity of group communication cl disruptions. By harnessing the power of decentral d tech ologie and cryptographic primitives, dynamic group ust, scal membership management protocols can provide r e, and secure solutions for modern distributed systems.

In the rapidly evolving landscape of digital communication, ensuring the security and scalability of group membership management protocols is paramount. As organizations incl singly rely on collaborative environments and distributed res in group membership while maintaining robust security measures systems, the ability to manage dynamic ch s the d becomes essential [7] [8]. This necessitat lopment of innovative protocols and technologies that can address the challenges posed by dynamic gr aures nd evolving security threats. One of the fundamental aspects of group communication protocols is at of secure communication channels among multiple parties [20] e establich [18]. Traditionally, cryptographic procols such as group key distribution schemes have been employed to facilitate  $[\mathbf{h}]$ secure communication wit [20]. However, existing protocols often face limitations in scalability and n gro security when confronted ith dyna ic changes in group membership, such as the addition or removal of members. These limitations c uine th confidentiality, integrity, and availability of group communication channels, und es to me security of sensitive information and organizational operations. posing signi *c*halle

To challenges, researchers and practitioners have explored novel approaches to dynamic group member nage, ot, leveraging advancements in cryptography, distributed systems, and blockchain technology. the integration of CGKA for group formation, which revolutionizes the collaborative process of One su oprð up keys. By ensuring that each member actively contributes to the generation of the group key, CGKA nerating d collaboration within the group, enhancing the security and resilience of the group's cryptographic fo trust infrast e. Additionally, the employment of LDH for key generation offers a robust and efficient method for w deriving shared secret keys. LDH leverages the mathematical properties of lattices to generate keys, ensuring confidentiality and integrity in communication channels. This approach enhances the security of group communication protocols by providing a secure foundation for key generation, even in the presence of dynamic changes in group membership.

The integration of blockchain technology for implementing membership changes introduces a decentralized and transparent approach to managing group dynamics within distributed systems. By leveraging blockchain's distributed

ledger technology and smart contracts, membership changes can be executed securely, transparently, and efficiently, enhancing the integrity and resilience of the group's membership management system. This enables secure addition and removal of members from the group, ensuring that the cryptographic infrastructure remains robust and scalable in dynamic group environments.

The key contributions of the article is,

- The integration of CGKA transforms the process of generating group keys by ensuring a live participation from all members. This collaborative approach instills trust and fosters collaboration aroung group members, enhancing the overall security and resilience of the group's cryptographic in the constant of the cons
- Leveraging LDH for key generation provides a robust and efficient method for securely driving sheed secret keys. LDH utilizes the mathematical properties of lattices to ensure confident lity and intentry in communication channels, thereby strengthening the security of the group's cryptur raphic of the arous.
- The incorporation of blockchain technology introduces a decentralizer and tansport approach to managing group membership dynamics. By leveraging blockchain's distribute dedger to hnology and smart contracts, membership changes can be executed securely, transport and efficiently.
- The integration of CGKA, LDH, and blockchain technology presents a comprehensive solution for advancing the security and scalability of dynamic group membership management protocols. This comprehensive solution offers a robust framework for secure and efficient communication in contemporary environments, addressing the complex maller is associated with dynamic group membership management.

The remainder of the article includes related works, mobile attement, includology and results in section 2, 3, 4 and 5. The paper is concluded in section 6.

#### 2. Related Works

The original purpose of VANETs, was to help with fic control and safety communication [21]. Owing to the notable advancements in contemporary automobiles, VACT functions have broadened to encompass pertinent services related to infotainment and comf c. The necessity to safeguard them has grown even more as a result of this up key is essential to VANET protection. In extremely volatile growth. Transparent sharing of a crypterar systems like VANETs, it is challe ging to revis the group key on a regular basis due to the rapid changes in difficult o create a group management key mechanism that is safe, scalable, membership in a group. It is there and effective. The high pro ases associated with group key computing and extraction, extra processing ex and interpersonal overhea when g up availation changes, as well as receiving vehicle collaboration are only a few of the restrictions introdu nt GKM methods. This study presents a unique GKM mechanism, ALMS, to l by cui solve these vestigation shows that because ALMS involves a minimal computational cost for the entire TA the per who receives vehicles, it is quicker to implement than current protocols. Furthermore, it is not co the key distribution issue that symmetrical key management techniques are. Furthermore, the inec only burder adds to the TA for affiliation changes is a little one. This is accomplished by separating the t AL from the group key calculation and carrying it out offline so as to preserve the encryption group initializ on pr cret's siz

D communications is a prospective 5G method for dynamic situations that can enhance the effectiveness of ading messages for group interaction because of the adaptability of devices [22]. Furthermore, each service on an ad no. hetwork is a programme running on the VANET. In order to minimize latency throughout vehicle discussions, communication between vehicles is being implemented in ad hoc contexts, including IoDs networks, C-V2X modules, drone fleet supervision, and autonomous driving systems. Nonetheless, facilitating safe and efficient interaction among teams is a pressing issue. It suggest a decentralized ledger-based dynamic group administration tool as the answer to these issues. The research presented here shows that a structure that is hierarchical built around distributed ledgers can handle dynamic groups more quickly and easily, without sacrificing security and functionality. Moreover,

the suggested approach can lessen the possibility of a single point of failures by facilitating the flow of data via immediate interaction without the need for a centralized database. Furthermore, an outside organization with deep ties to Taiwan's top automobile electronic suppliers globally conducted testing on the findings.

New approaches to access control have emerged in response to the explosive proliferation of IoT devices handling private information, with the goal of protecting this information from unauthorized usage [23]. To guarantee safe data delivery to authorized users, a dynamic Internet of things context that is marked by a high signaling overhead d users' movement poses a serious challenge. Therefore, GKM serves as the essential method for controllin the assignment of keys for controlled access and safe sharing of data during these dynamic contexts. Unfortunately majority of the GKM-based access management techniques now in use for the IoT depend on centra making them unable to handle the scalability issue brought on by the huge amount of IoT devices and g wing nui of subscribers. Furthermore, neither of the GKM methods in use today encourage group members adiv ey just concentrate on dependent asymmetric group keys to communicate inside each subg them in ective for subscribers exhibiting extremely dynamic behavior. It provide a unique DLGK C to dress se issues. The suggested system improves the administration of users' subgroups and reduces f retypin burden of the KDC by using a hierarchical design made up of many SKDCs and one KDC. Additional rand-new master's token management technique is presented to control the distribution of keys among users Vith this type of protocol, join/leave events have less overhead in terms of processing, storing, and transmissio ering the strain brought By N on by reentering at the core system, the suggested method allows for e Internet of Things design and counteracts the risk of only having one point of failure.

Resolving the speed constraint of PoW-based blockchain vpically enable just hundreds of approval, is a major goal of the PBFT agreement transactions per second and take moments to months f ach method [24]. PBFT is generally used in tiny network becau of its or node expansion caused by numerous internode connections. In this paper, an extensible m layer J FT-based consensus process is suggested for enabling PBFT in big structures, like blockchain and enor oT ecosystems. The technique works by systematically aggregating nodes into distinct levels and restricting transition inside the group. First provide an ideal double-layer PBFT and demonstrate a considerable reduction of communication difficulty. In particular, researchers demonstrate that interaction difficulty is minimized wided the nodes are spread equally throughout each sub-group in the following layer. FPD and FND metho accordingly, to analyses the safety threshold. In addition, researchers offer a workable procedure for the ggested double-layer PBFT systems. Lastly, the findings are and containications efficiency in arbitrary-layer PBFT platforms. The efficacy expanded to include security analysis of the analytical data is cont by results of simulations.

curity and scalability in various technical contexts are covered in the literature Three main topics that mprove ip in VANETs is dynamic, the study emphasizes the difficulties in safely study. First off. maintaining bgraph group keys. It suggests a novel GKM technique known as ALMS, which solves scaling problem s computing expenses. Second, the emphasis moves to D2D communications in 5G contexts, highlighti . ity of effective and safe group collaboration. A decentralized ledger-based dynamic group e nec s used in the suggested approach to increase agility without sacrificing security. The study manag ent B the a discussion of control of access in Internet of Things contexts, where scaling and uniqueness issues concludes centralized GKM approaches. It presents a DLGKM-AC systems that lowers processing complexity for typ ar ing subgroup administration and adaptability. Every research offers fresh strategies to tackle certain while blems, advancing both safety and scalability in the fields they study.

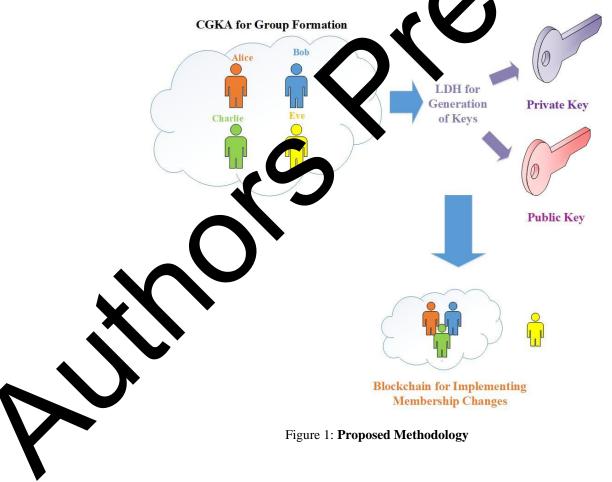
#### 3. Problem Statement

In contemporary digital environments, managing dynamic group membership while ensuring robust security measures poses significant challenges. Existing group membership management protocols often struggle to adapt to dynamic changes in group composition, leading to scalability issues and security vulnerabilities. The need for secure

and scalable group communication channels is paramount for organizations operating in collaborative environments. Therefore, there is a pressing need to develop innovative protocols that can effectively address the complexities of dynamic group membership management while maintaining the security and integrity of communication channels [25]. The proposed methodology aims to address these challenges by integrating CGKA for group formation, leveraging LDH for key generation, and employing blockchain technology for implementing membership changes. This comprehensive approach seeks to revolutionize group membership management by fostering trust, enhancing security, and ensuring scalability in dynamic group environments.

### 4. Proposed Dynamic Group Membership Management

Integrating CGKA for group formation ensures that each member actively contributes to the ge ration of he group key, fostering trust and collaboration within the group. This approach enhances the securi silieng of an g all p the group's cryptographic infrastructure by distributing the responsibility of key generation ants. af LDH is employed for the generation of keys, leveraging the mathematical propert ces ecurely derive ot la shared secret keys. LDH provides a robust and efficient method for generating ographic applications, vs in cry ensuring confidentiality and integrity in communication channels. Employ ekchain technology for implementing membership changes offers a decentralized and transparent approach, and ving for the secure addition and removal of members from the group. By leveraging blockchain's distribute led technology and smart contracts, membership changes can be executed securely, transparently, and effic itly, enhancing the integrity and resilience of the group's membership management system. It is depicte



#### 4.1 Integrating Contributory Group Key Agreement for Group Formation

The role of Contributory Group Key Agreement (CGKA) in group formation is paramount for establishing robust and secure communication channels among multiple parties. CGKA facilitates the collaborative generation of a group key, ensuring that each member actively contributes to the process. This collaborative approach enhances the overall security of the group by distributing the responsibility of key generation among all participants, mitigating the risk of a single point of failure. By involving each member in the key generation process, CGKA fosters a sense of trus and accountability within the group, as every member plays a crucial role in establishing secure communication changels.

CGKA enables dynamic group membership management, allowing new members to join or exi ng me ers to leave the group without compromising the security of the group key. This flexibility is essent adaptin to changing group dynamics and ensures that the group key remains secure even as the composition the g ves over time. Additionally, CGKA provides resilience against attacks and unauthorized a as the group key mp is derived collaboratively from the individual contributions of all members. Our all, the KA in group le of formation is to facilitate the collaborative generation of a secure group key, i ust, accountability, and ring resilience within the group.

The Alice, Bob, Charlie, and Eve group formation mimics a situation similar quantum group key distribution (QKD), which is a crucial idea in quantum cryptography that aims to proe channels of communication. In this configuration, Bob and Charlie separately produce random bases, an rates random secret data bits and Ali encodes them using a randomly selected basis. An important when Eve, the eavesdropper, isk a intercepts and perhaps modifies the qubits that Alice Charlie in order to create a secure transn to Bo communication channel. This hypothetical situation gnificance of secure communication protocols .pha s th in quantum cryptography and the continuous str ide robet defenses for confidential data in quantum gle to pr communication networks.

#### 4.1.1 Alice

Alice has a crucial part in the est dishment of the group since she is the one who starts and sends the quantum communication process. Secure com a is tablished when Alice creates random secret data bits and encrypts them using a randomly selected sis. Alice process use of quantum mechanics to increase the security of the communication channel by encodi the secret bits in quantum states. Her proactive engagement guarantees the integrity and secrecy of the ansferred, laving the groundwork for Bob and Charlie to safely extract the ion secret data bits by receiving and ded ling the quantum states. Alice's role emphasizes how important it was for her to build safe communication ghlighting how important it is for quantum cryptography to have strong security nnels mechanism

#### 4.1.2 B.

Bob phase servicial part in the establishment of the group as he is one of the intended recipients of the quantum communication that Alice started. Bob is essential to Alice's process of receiving and decoding the quantum states that he send the order to get the bits of secret data. Furthermore, Bob creates random bases on his own, strengthening the computation channel's security even further. The unpredictable and complicated nature of these independently unpredictable bases makes it more difficult for possible eavesdroppers, like Eve, to collect and decode the sent data. Bob's enthusiastic involvement underlines how secure communication protocols are collaborative in nature, emphasizing how crucial it is for numerous parties to cooperate in order to create and maintain safe channels in quantum cryptography.

#### 4.1.3 Charlie

Charlie plays a critical part in the establishment of the group as an additional intended recipient of the quantum communication that Alice started. Charlie, like Bob, has to receive the quantum states that Alice sends and decode them in order to get the bits of hidden data. To add to the variety and unpredictable nature of the communication process, Charlie also independently creates random bases. By adding more randomness and complexity to the communication channel, this independent base generation strengthens its security by increasing the difficult for possible eavesdroppers, like Eve, to collect and decode the sent data. Charlie's active participation highlight the cooperative aspect of secure communication protocols, emphasizing how crucial group efforts are to a the preserving secure channels in quantum cryptography.

## 4.1.4 Eve

Eve plays the role of the eavesdropper in the group formation, which pres y risk to the 18 quantum communication process that Alice started. Eve's main goal is to interce be modify the quantum ind m states that Alice sends to Bob or Charlie in an effort to obtain the secret data bits with ing discovered. Eve listens in on the communication channel with the intention of using system flaws and vulnerability to get private information and jeopardize the communication's integrity. Eve's existence highlights the pers culties in guaranteeing nt di. e significance of strong security secure communication in quantum communication networks and emp mechanisms in quantum encryption to identify and neutralize eavesdry efficiently. bin

# 4.2 Lattice Diffie–Hellman for Generation of Keys

LDH is a cryptographic protocol utilized for eratin keys urely, leveraging the mathematical properties of lattices. The protocol begins with each party, pically ferred to as Alice and Bob, independently generating random matrices and vectors. Alice generates a rai matrix and a secret vector, while Bob generates another random matrix. Alice computes a noisy vector by adding undom noise to the result of a matrix-vector multiplication, and she sends this noisy vector to Bob. Upon receiving a noisy vector, Bob computes another noisy vector by multiplying it with his random matrix and agoing more random noise. Bob then sends this noisy vector back to Alice. Finally, Alice can compute the shared se et k erforming an inner product operation between the received noisy vector and her secret vector.

ed secret key is securely generated over an insecure communication channel This process ensures that the si without directly exchanging ate formation. The security of the LDH protocol relies on the hardness of the any LWE problem, which ma s it com tationally infeasible for an eavesdropper to recover the shared secret key from the exchanged nois By le raging the mathematical properties of lattices and the difficulty of solving the cto LWE proble protocol provides a robust and efficient method for generating keys in cryptographic s LD fidentiality and integrity of communication channels. Table 1 shows the parameters of applicat ng the c LDH.

	Parameter	Description	
	Lat e.Br is	Represents the mathematical structure of the lattice, typically defined by a basis matrix.	
	Secret vector	Random vector chosen from a discrete Gaussian distribution, used to compute the public	
		key.	
	Public Key	Vector obtained by taking the inner product of the lattice basis vectors with the secret	
vector.		vector.	
Noise Term Small noise term added to the in		Small noise term added to the inner product computation to ensure the resulting key is	
•		indistinguishable from random.	

#### Table 1: Parameters of LDH

Encryption Scheme	Utilizes the computed public key and the recipient's private key to encrypt messages	
	securely.	
Decryption Scheme	Uses the recipient's private key and the sender's public key to decrypt encrypted	
	messages.	

# 4.2.1 Generation of Private Key

secret

The generation of a private key is a fundamental aspect of asymmetric cryptography, where each party communication session possesses a unique key pair consisting of a private key and a corresponding public key private key is a securely generated, random string of binary digits or alphanumeric characters, typic using cryptographic algorithms and protocols. The process begins with the selection of a secure ra dom nui generator (RNG), which ensures that the private key is generated with sufficient entropy to ptogi hic sist attacks. The private key is then generated by the RNG and stored securely in the pos key houser. It is ol crucial that the private key remains confidential and is not shared with any unauthori in the security partie of the cryptographic system.

In asymmetric cryptography, the private key is kept secret and is known only to the owner, whereas the corresponding public key is shared with other parties for encryption or signature verification purposes. The private key plays a vital role in cryptographic operations such as decryption, digital signature generation, and key agreement protocols. Overall, the generation of a private key is a critical step in stablic ang secure communication channels, digital signatures, and other cryptographic operations, ensuring the considerability integrity, and authenticity of data in modern cryptographic systems.

Using the mathematical features of lattices, the LDH cuptographic protocol generates private keys in a secure manner. The Learning with Errors (LWE) issue, which argumentati is computationally difficult to retrieve a concealed secret from a given set of noisy linear equations, is uncoundation of the protocol. Using the use of an unsecure communication channel, two people, known as Alice and Bob, want to construct a shared secret key using the LDH protocol.

Alice creates a secret vector (s) and a recommutrix (A) in the first phase of the LDH procedure. The random matrix A is a n x m matrix with every randomly elected elements from a vast field. The private key is contained in the m-dimensional secret vector, a Vs. After mat, Alice multiplies the matrix-vector result by a little amount of random noise to get a noisy  $e_{-x}$ . In terms of math, this is expressed as:

$$e = As + noise \tag{1}$$

In the susception process, Bob receives the noisy vector e from Alice over the unsecure communication channel. After geing d = 0 Bob craces a new random matrix B with dimensions  $d=F \times mn \times m$ . He then multiplies d = e by dB and action new stor frandom noise to create a second noisy vector, d = f. In terms of math, this is expressed as:

$$f = Be + noise \tag{2}$$

It last bob uses the unreliable channel to give Alice the noisy vector f back. Alice may then obtain the shared by calculating the inner product of the vector f and her secret vector s after obtaining f. In terms of math, expressed as:

$$s_{shared} = f. s \tag{3}$$

Because the LWE issue is hard, it is computationally impossible for an eavesdropper to get the shared secret key *s* from the noisy vectors *i* and *f*. This is the foundation for the security of the LDH protocol. The LDH protocol offers

a safe and effective way to generate private keys for use in cryptographic applications by taking use of the mathematical characteristics of lattices and the challenge of addressing the LWE problem.

#### 4.2.2 Generation of Public Key

The generation of a public key is a fundamental process in asymmetric cryptography, where each participant in a cryptographic system possesses a unique key pair consisting of a public key and a corresponding private key. Unlittle private key, which must be kept secret, the public key is intended for distribution and is made freely available to other parties. The generation of a public key typically involves applying mathematical algorithms and protocor to derive a value that is mathematically related to the corresponding private key. This relationship ensures that mesage encrypted with the public key can only be decrypted by the corresponding private key and vice verse provides a mechanism for secure communication and digital signatures.

One of the most common algorithms used for public key generation is the Bort subprise which involves selecting two large prime numbers and performing mathematical operations to generate a public key openent and a corresponding private key exponent. The public key consists of the modulus and the public apponent, while the private key consists of the modulus and the private exponent. Overall, the generation of a public key is a crucial step in establishing secure communication channels, digital signatures, and other cryptographic operations, ensuring the confidentiality, integrity, and authenticity of data in modern cryptographic systems.

LDH is a cryptographic protocol used for generating public keys ocurely based on the mathematical properties of lattices. The protocol leverages the hardness of the LWE problem who be ates that it is computationally difficult to recover a hidden secret from a given set of noisy linear equation on LDL upon arties, typically referred to as Alice and Bob, aim to establish a shared public key over an interest communication channel.

In the first step of the LDH protocol, Alice generates a roldom matrix B and a secret vector s. The random matrix B is a n x m matrix with elements chosen uniformly at a roldom from a large field. The secret vector s is a m-dimensional vector containing the private key. Alice then computes a new vector d by adding a small random noise to the result of the matrix-vector multiplication *Bs*. Mathematically, this can be represented as:

$$Bs + noise \tag{4}$$

Next, Alice sends the noisy entor d to Bel over the insecure communication channel. Upon receiving d, Bob generates another random matrix C on limensions n x m and computes a second noisy vector f by multiplying d with C and adding another set of random noise Wathematically, this can be represented as:

$$f = Ce + noise \tag{5}$$

Bob then as ds the being vector f back to Alice over the insecure channel. Upon receiving f, Alice computes the inner product of the vector f and her secret vector s, resulting in the shared public key  $R_{shared}$  mathematically, this can be represented as:

$$R_{shared} = f. s \tag{6}$$

The structure of the LDH protocol relies on the hardness of the LWE problem, making it computationally feasible for an eavesdropper to recover the shared public key  $R_{shared}$  from the noisy vectors d and f. By leveraging it is matter and the difficulty of solving the LWE problem, the LDH protocol provides a secure and efficient method for generating public keys in cryptographic applications.

#### 4.3 Employing Blockchain for Implementing Membership Changes

Employing blockchain technology for implementing membership changes offers a decentralized and transparent approach to managing group dynamics within distributed systems. Blockchain, as a distributed ledger technology, maintains a tamper-resistant record of transactions across a network of nodes. Each transaction, including membership changes such as additions or removals of members, is cryptographically signed and recorded on the blockchain, ensuring transparency and immutability. When a new member seeks to join the group, a transaction is created and broadcasted to the network, detailing the necessary information for membership approval. Similarly, when a member needs to be removed from the group, a corresponding transaction is generated, reflecting the change in member ip status.

Blockchain smart contracts can automate the process of membership changes, executing d rule logic to validate and authorize membership requests. Smart contracts can enforce member verif ities and ensure compliance with predefined rules before processing membership cha blockchain's tion decentralized nature eliminates the need for a central authority to manage member g the risk of nip chap s. reduc chnology, implementing single points of failure and enhancing the resilience of the system. By leveraging bloc membership changes becomes more transparent, auditable, and secure, providing a ro t framework for managing group dynamics within distributed systems.

#### 4.3.1 Display Keys Before Eve is Removed

Before Eve is removed, the display of keys showcases c contributions of each participant ograp in the group, including Alice, Bob, Charlie, and Eve, Each nber's oution to the group key is visually represented, illustrating their individual role in the k ocess. The display highlights the collaborative tio nature of the key generation scheme, emphasizin each member's contribution in ensuring the the imp tance lay serves as a visual aid for monitoring and verifying security and integrity of the group key. Additional the d the distribution of cryptographic responsibilities with group, providing transparency and accountability in the key generation process.

#### 4.3.2 Remove Eve from the Grou

Using blockchain techno Eve from the group involves executing a series of transactions on move the blockchain network to upda he group's embership records and revoke Eve's access privileges. First, a transaction is created to initiate the noval process, specifying Eve's identification details and the reason for her removal. This transaction i the blockchain network, where it is verified and added to the blockchain's immutable ledger by th consensus mechanism. Smart contracts deployed on the blockchain can networ s and logic to validate the removal request, ensuring that it complies with the automatically execute pre fined ru group's men cedures. bersh

Once he remote transaction is confirmed and added to the blockchain, Eve's access privileges are revoked, as her contographic contributions to the group key are invalidated. This ensures that Eve no longer has access to the group's resources or confidential information. The removal process is transparent and auditable, allowing all memory of bace out to verify the transaction and confirm Eve's removal from the group. By leveraging blockchain technology the removal of Eve from the group is executed in a secure, transparent, and decentralized manner, entracing threategrity and resilience of the group's membership management system.

# **3 Public Keys before Eve is Removed**

Before Eve is removed, the display of public keys showcases the cryptographic contributions of each part icipant in the group, including Alice, Bob, Charlie, and Eve. Each member's public key is visually represented, illust rating their individual role in the cryptographic operations within the group. The public keys serve as essential comp onents for encrypting and decrypting messages, establishing secure communication channels, and verifying digital si gnatures. The display highlights the collaborative nature of the group's cryptographic infrastructure, emphasizing the importance of each member's contribution in ensuring the security and integrity of the group's communication protoc ols. Additionally, the display serves as a visual aid for monitoring and verifying the distribution of cryptographic res ponsibilities within the group, providing transparency and accountability in the cryptographic operations.

#### 4.3.4 Public Keys after Eve is Removed

After Eve is removed from the group, the display of public keys reflects the und contributions of the remaining participants, namely Alice, Bob, and Charlie. With ey removed, the ıhİ display now showcases the public keys of the remaining members, illustrating the gement in the contin ed in group's cryptographic operations. The removal of Eve ensures that only trusted ontribute to the group's mbers cryptographic infrastructure, enhancing the security and integrity of the communication innels. The updated display serves as a visual confirmation of Eve's removal from the group and reinforces the colla tive nature of the group's cryptographic protocols. Additionally, it provides transparency and accountability ptographic operations, allowing all members to verify the distribution of cryptographic respons hin the group.

$$Bob's Updefour white field = [4\ 6\ 1] \tag{11}$$

Charties Update Public Rey = 
$$[2 \ 8 \ 2]$$
 (12)

#### Algorithm: Dynamic Group Membership Mangement

#### Initialize

Set the group size to n. Generate a random prime number Choose a generator g for the cycle group Zp

#### **Key Generation**

Each member generates of the second second second as sk[i], where i represents the member's index in the group Calculate the corresponding public keys or each member:

#### **Group Key Agreement**

Each member be clease their ablic key pk[i] to all other members in the group.

Upon receiving all public keys, each member computes the group key as follows:

#### LDH Key Generation

*Choose fice progeneters and generate a lattice basis.* 

Each memory generates a random vector s[i] as their secret key.

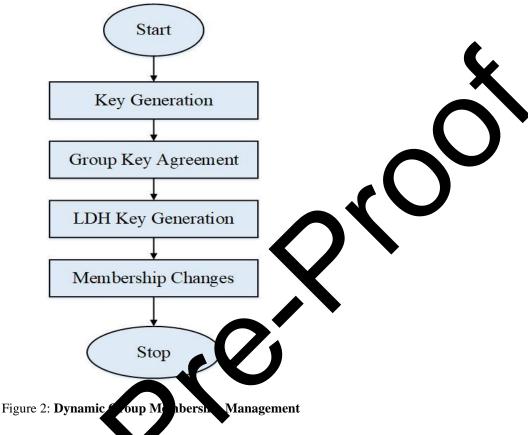
- Compute the presponding public key for each member using LDH
- Share to public keys with all other members.

#### nbership Changes with Blockchain

e a shart contract on the blockchain to handle membership changes.

When the member joins, they submit a transaction to the smart contract

*A a member leaves the group, they submit a transaction to revoke their membership Membership changes are recorded on the blockchain, providing transparency and accountability* 



#### 5. Results and Discussion

The proposed method is implemented in Python soft and the efficiency is evaluated and compared with existing protocols. The performance evaluated is given in this section.

# 5.1 Network Graph of Participation with Keys

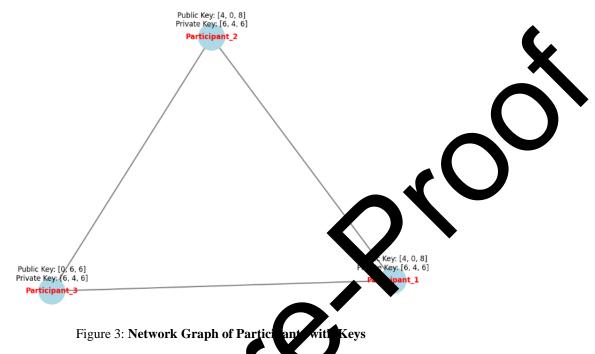
A network graph of participant with keys illustrates the relationships between participants in a cryptographic system, showcasing their respects put is and private keys. It provides a visual representation of the key distribution within the network, facilitating an usis of key sharing and ensuring the integrity and security of cryptographic communications.

Partic onts	Public Key	Private Key	
Naticipal. 1	0,6,6	6,4,6	
Part. Pant 2	4,0,8	6,4,6	
rticipant 3	4,0,8	6,4,6	

<u> </u>	-		
Toble 2. Notwork	Chank	of Doutioinon	to with Vora
Table 2: Network	(Trann	ог вагисирай	is while Kevs

Table 2 depicts the network graph of participants along with their corresponding public and private keys. The public keys, representing the shared information accessible to all participants, are listed alongside the private keys, which are kept confidential and unique to each participant. Observing the network graph, it becomes evident that Participants 2 and 3 share identical public and private key pairs, suggesting a potential redundancy or oversight in the key generation process. This uniformity may raise concerns regarding the uniqueness and security of the cryptographic keys within the network, warranting further investigation into the key generation methodology and ensuring the integrity of the cryptographic framework. It is depicted in Figure 3.

#### Network Graph of Participants with Keys

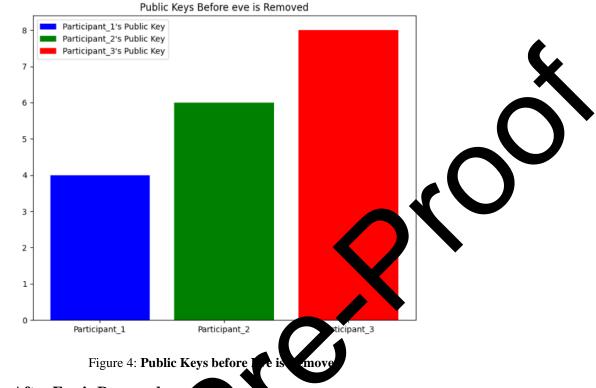


#### 5.2 Public Keys before Eve is Removed

The public keys before eve is removed gr in Figu 4 disp. s the public keys associated with Participants 1, 2, and 3 along the x-axis. Each participant's put key represented by on the graph. The y-axis represents the e group, the graph shows the distribution of public keys value of the public keys. Before Eve's removal from generated by each participant, reflecting their contribution the group key agreement process. Analyzing this graph allows for visualizing the diversity and distribution of public keys across participants, providing insights into the key. Additionally, it facilitates monitoring any irregularities or cryptographic strength and security of le group anomalies in the public key distribut indicate unauthorized access or compromised participants within the group.

Tat 3: Public Keys before Eve is Removed	
Pa ticipants	Public Key
Alice	4
Bob	6
Charlie	8

the participants' public keys prior to Eve being kicked out of the group. The numbers 4, 6, and **`ab**I Alice. b, and Charlie's respective public keys. These public keys are crucial for creating safe channels 8 repres tion within the organization. Each member of the group has contributed differently to the group's ommun as seen by the variances in the size of their public keys, which signify their distinct responsibilities cry ithin cryptographic infrastructure. The public keys are essential for message encryption, safe connection ment, and digital signature verification. This emphasises the need of each member's participation in maintaining the security and integrity of the group's communication protocols.



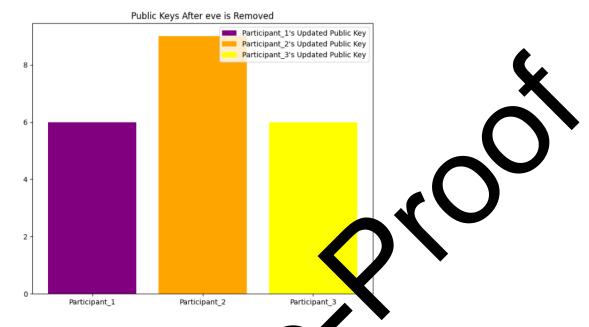
# 5.3 Public Keys After Eve is Removed

The public keys after eve is removed grappin Figure 5 depicts the public keys associated with Participants 1, 2, and 3 on the x-axis. Each participant's public keys prepresented by a distinct line or data point. Following Eve's removal from the group, the graph illustrates the update distribution of public keys generated by the remaining participants. This visual representation allows for observing any changes or adjustments in the distribution of public keys after the removal of a compromiser or unauthorized participant. Analyzing this graph facilitates assessing the impact of Eve's removal on the security adjuteger of the group key agreement process, providing insights into the resilience of the group against potential security the eats or attacks.

#### e 4: Public Keys after Eve is Removed

Palicipants	Public Key	
, ice	4	
Bob	6	
Carlie	8	

The public keys of the participants in Table 4 do not alter following Eve's expulsion from the group. The public keys that the, Bob, and Charlie still have are denoted by the numbers 4, 6, and 8, respectively. The group's coppograph infrastructure is stable and intact after Eve was removed, as indicated by the consistency of the public keys. The surviving members of the group continue to provide cryptography in the same way, thus safe lines of communication continue even in the event of a shift in group dynamics. This highlights the robustness of the group's integraphic protocols and the efficiency of the systems set up to handle membership changes in an open and safe mannet.



#### Figure 5: Public Keys after Eve is Pane

In terms of key size and computational cost, Table 5 g gested LDH approach with other the s methods that are already in use. With a 256-bit key and a con Nexity of 10<sup>5</sup> operations, the LDH tationa technique shows competitive performance, indicating and applicability for cryptography applications. iene ECDH have bigger key sizes (2048 and 256 On the other hand, conventional techniques like RS Key Ex hange bits, respectively) and have computational complex of 10<sup>6</sup>  $0^7$  operations. With a key size of 512 bits, PQC likewise offers a competitive option; nevertheless, its compute demands are comparable to those of RSA Key Exchange. The comparison highlights the importance of LDH as jable method for safe key exchange, providing a balance between computational performance, key size, and

Table .	Comparison	with	Existing	Methods
			-	

Method	Ke Size (bits)	Computational Complexity (Operations)
FODU	256	106
RSA key Extrange	2048	107
PQC	512	107
Prop. 1. OH	256	105

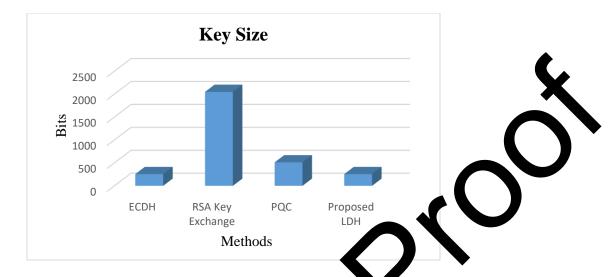


Figure 6: Comparison with Existing Methods

## **5.4 Discussion**

The results of the comparison between LDH and existing met core the efficacy of LDH in terms of key size and computational complexity. LDH demonstrates co perf mance with a relatively small key size of 256 bits, comparable to other contemporary cryp ethods. is compact key size is advantageous for various applications, including resource-constrained re efficient utilization of computing resources aviron ents is paramount. Additionally, LDH exhibits a low compu ional complexity, with an order of magnitude fewer operations required compared to traditional methods A Key Exchange and ECDH. This reduced computational overhead makes LDH particularly appealing for scenario where computational efficiency is critical, such as real-time communication systems or high-throughput data processing vironments.

The results suggest that LDH nising solution for addressing the challenges posed by emerging quantum computing threats. By lev ematical properties of lattices and the hardness of lattice-based ne ma problems, LDH provides a robust nework secure key exchange, even in the presence of quantum adversaries. The relatively small key size imputational complexity of LDH further contribute to its suitability for postlow quantum cryptographic erall, the results highlight LDH as a viable alternative to traditional blicati mpelling combination of security, efficiency, and resilience against quantum cryptographic methods, o ering a d computing a tack

### 6. Conclusion and Future Works

a comprehensive solution for addressing the challenges of managing group membership resen inporary communication systems. By integrating CGKA, LDH, and blockchain technology, the dvnami in d col extension enhances the security, scalability, and efficiency of dynamic group membership osed n protocols. CGKA ensures that each member actively contributes to the generation of the group key, mar ast and collaboration within the group. LDH provides a robust and efficient method for generating keys, sterin ing the mathematical properties of lattices to ensure confidentiality and integrity in communication channels. Additionally, blockchain technology offers a decentralized and transparent approach for implementing membership changes, allowing for secure additions and removals of members from the group while maintaining the integrity of the cryptographic infrastructure. The future research could focus on further optimizing and refining the proposed protocol extension to enhance its performance and effectiveness in real-world scenarios. This could involve conducting more extensive simulations and performance evaluations to validate the scalability and efficiency of the protocol extension under various conditions. Additionally, research could explore the integration of advanced cryptographic techniques and protocols to further strengthen the security of dynamic group membership management. Furthermore, investigating the impact of emerging technologies such as quantum computing on the security of the protocol extension could provide valuable insights into potential vulnerabilities and mitigation strategies. Overall, continued research in this area is essential for advancing the state-of-the-art in dynamic group membership management protocols and ensuring the security and scalability of communication systems in dynamic environments.

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