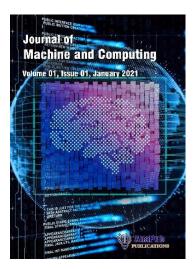
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Implementation and Enhancement of Endogenous Security Mechanisms in Photovoltaic Data Storage and Transmission

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Abstract—The global energy market is mi and to rd sustainable renewable energy sources (RES), with solar energy (SE) being the most significant due to its abundance and reliability. Photovoltaic (PV) converts SE into electricity relying on data integrity and security. However, digitized data has cybersecurity vulnerabilities, reluding data breaches and attacks. Traditional security systems can provide essertial tection but fail to address PV's dynamic and distributed nature, leading to gaps in effense against evolving cyber threats. The study proposes an endogenous security model for improving data transmission and storage within PV. It uses a Verification Feedbace Mechanism (VFM) to integrate routing methods, compute efficient data ric scally, and verify their integrity. The model also incorporates m paths, schedule ev infinit ructure and key management protocols to ensure secure data transmission cryptographic and mana, ment. This approach addresses challenges in data forging and ensures the integrity of rk's **U**nponents. The study compared two methods and found one model superior in the net communication integrity and system adaptability. It achieved latency statistics below 20 ms and gintamed Network Throughput (NT) at 9.2 Gbps even when attacked, demonstrating its effectiveness in securing PV from multiple cyberattacks.

Keywords—Photovoltaic Systems, Renewable Energy Sources, Solar Energy, Cyber-Physical Security, Data Transmission Integrity Rates.

I. INTRODUCTION

Since the global community is migrating towards Renewable Energy Sources (RES), photovoltaics (PV) has evolved into a significant component of the Renewable Energy (RE) geographical region. By reducing the requirement to rely on natural resources, PV automatically transfers energy from ultraviolet radiation into electrical power, resulting in a more sustainable RE. Solar PV [1] differs among RE uses in that devices can be built up or down to connect to be Smart Grid (SG) at multiple levels, from individual residences to enormous solar power platts. Among other key measurements, data on PV radiation, electricity generation, and EC to be is necessary for PV functioning, management, and efficiency improvement [2] The nata improves the reliability and sustainability of PV installations, enabling provide maintenance of these systems via proper data processing. In order to successfully monitor and mantain PV, it is essential to maintain the reliability and safety of these vital data.

existing network connectivity However, digitizing and transmitting such data throu h th infrastructure to perform all the PV-related operations has expected the system to a complex e thruts to PV can range from data breaches and environment of cybersecurity attacks [3]. unauthorized access to more complex a. cks promising the integrity and availability of critical PV energy data [4]. Further, as these F are ultimately integrated into the national Smart Grid (SG), the complexity of handling and the probable impact of these attacks have become more imminent and have extended beyond the lividual installations, which posed risks to the broader energy infrastructure's stability and religibility. In response to these challenges, various security mechanisms have beer lop I based on techniques such as encryption, authentication, and secure communicatio protocos. Such security mechanisms only frequently addressed the specific itv had provided room to be exploited by knowledgeable attackers [5]. aspects of zbe sec

This is when endogenous security systems come into the field, which is unlike other security measures had are applied as external layers of security; the endogenous security systems, on the other had are models that are integrated into the core operational model of PV [6] It is implemented as core components of the system's design thereby ensuring the security and data therein the model during the data transmission and storage processes. However, few models have been developed using endogenous principles for PV data security, which provides plenty of room for proposing models in this domain. The motivation for this work is grounded in the

limitations of existing security systems and the limited work on endogenous-based cyber-physical systems (CPS) for PV environments.

The proposed work addresses the above limitations and gaps by introducing an endogenous security model for the security and integrity of data transmission and storage in PV. The model employs an integrated routing model that employs three routing strategies: Data Integrated ity Forwarding (DIF), Load Balanced Forwarding (LBF), and Path Diversity Forwarding (\mathbf{PI}) together with Verification Feedback Mechanism (VFM) for ensuring reliability. The integr ed model enables the simulation to take advantage of the best possible use of the esour network while reducing the probability of data manipulation [7]. In aluating the .dditi to validity by finding paths with minimal collision and actively update th e paths according to network situations, the design includes techniques for computing the most appropriate data paths for scheduling [8]. The design employs a secret cryptographic key estem and protocols for key management to ensure security when transmitting information, a multiple routes. Based on the findings of the test of the security system using different performance indicators, the recommended scores ranging from 97% to 99%, reliability security model attained data transmission reabili Aghput (NT) maintaining near 9.2 GBPS under rates ranging from 94% to 96%, and Netw Th attack issues with latency less than 20 ms.

The article is organized as follower Section 2 provides the existing literature review, Section 3 shows the historical context of the weak Section 4 presents the recommended security model, Section 5 analyses the model and Section 6 presents the conclusion. The paper is structured in the following order.

II. LITERATURE REVIEW

Based to the enclored experienced by CPS-based PV, their [7] paper provided an in-depth review. The study has unphasized the diversity of cyberattacks that have been the target of PV, which have angeodrom data integrity to software-based attacks, and the work has also introduced a success rate metric to assess the impact of these attacks. It also explored model-based and datadrives approaches for threat detection and mitigation, highlighting the effective role of blockchain topology in securing software and CPS.

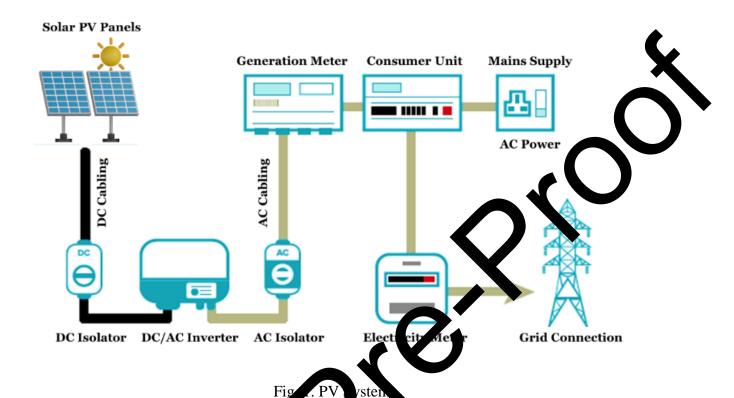
Authors [9] focused on the communication security of PV by a voltage regulation scheme. This scheme has been built to operate on two levels to reduce voltage deviation and voltage difference, and it incorporates a power compensation system for primary regulation and a consensus protocol

for secondary regulation. The approach was built to handle the communication topology changes and delays by proposing predictive compensation for packet loss and significant delays. Their method's effectiveness is being validated using MATLAB simulations, and the results have shown the models' better performance.

Authors [10] have addressed the broader implications of integrating RE sources into 5G, particularly the security problems associated with wireless data transmission and centralized pour trading. They indicate a secure energy market approach appropriate for Smart Grid (3G), which employs Wireless Sensor Networks (WSN) for communication and is endorsed by block bein.

In order to enhance the success of energy making and RE use, the a dual-chain nors opos design that saves electrical data about the blockchain while employing Sr t Contracts (SC) for business decisions. In Germany's "Digitalization of the Energiewende" germance [11] has dealt with the evolution of the distribution systems, including SG. The study seeks to show how bor e t Controllable Local Systems (CLS) and Smart Metres col build a secure energy data S F network, enabling secure communication with distributed h e PV cells and batteries. The reality that this reciprocal achievement meet data ecutive regulations in Germany highlights the chance for enhanced management and integration a decentralized RES into SG. Their study shows the development and execution of a node complete that deals with the secure communication requirements associated with cloud-based RES [12].

In order to verify the system second and reliable performance, the study emphasized the importance of security measures adapted to the demands of the network of communications. In an example experiment i ir model exhibited more significant enhancements in peakclipping, valley-fillin effects and overall load features, proving the value of the method. More ation of distributed energy and improving the effectiveness of RES, significan con. heer udying CPS difficulties associated with multi-station systems. The study includes [13] ha ŔЬ ource to supply a supplemental design and techniques for securing the natural several environment of Smart Energy Stations (SES). Security zone and congestion isolation systems are ues they suggest for improving SES-CPS in the context of potential attacks. two chni



r enhancing the quality and performance of Previous investigations have concentry ed . Photovoltaic Storage Systems (PVSS) in an Energy Blockchain (EB) context [14]. They achieve this by developing a task-matchin model employing the Genetic Algorithm – CLOUD-Gale-Shapley (GA-CLOUD-GS). The aev oped a method for integrating subjective and objective weights for matching tasks. everal emputer simulations, sensitivity studies, and comparison analyses have demonstrate the model's performance and provided helpful information for Publicly Verifiable Stret Sharing (PVSS) task balancing in an EB environment [15]. Applying SG to solar V then research focuses on threat identification and risk estimation. Their approach involve rend ring, assessing, and mitigating CPS risks specific to SG with those of solar PV integration. tilizing the proposed Scheme for Trans-disciplinary Research for India's Developing STRIDE) for threat classification and the following proposed model, the DREAD Fronomy Damage potential, Reproducibility, Exploitability, affected users, and Discoverability) stand for risk ranking model for prioritization, and the proposed study has been practical in the process of identifying the high-risk threats that have included information disclosure and elevation of privilege.

III. BACKGROUND

A. Structure of Solar PV

A typical solar PV (Fig. 1) comprises several key components that combine to convert sunlight into electrical energy that can be used in homes/industries in the power grid [16-20].

The components of the PV include:

- Solar PV Panels: The solar PV panels capture sunlight and convert it into Direct Current (panels) electricity. These panels contain PV cells made from semiconductor materials exhibiting effects.
- *DC Cabling and DC Isolator:* The DC electricity generated by the panels of transition via DC cabling. This wiring is connected to a DC isolator, a safety dence that disconnects the PV from the electrical circuit for maintenance.
- *DC/AC Inverter:* This device converts the DC electricity from the solar genels into Alternating Current (AC) electricity.
- *AC Isolator:* The AC isolator provides a point of disconnection for the AC converted from the solar panels.
- *Generation Meter:* The generation meter is conjected in the inverter and measures the amount of AC electricity the solar PV produces
- *Consumer Unit:* Also known as the fuse by the consumer unit is where the electricity is distributed to different circuits **constant** in the home.
- *Electricity Meter:* The electric content records the amount of electricity the SG consumes.
- Connection to the SG: The system is connected to the main supply SG.
- Mains Supply: The main upper represents the household's connection to the public electricity SG.
- B. Data Germution . Solar PV

Data cheraton in solar PV is a continuous process collected from the operators, owners, and utility computes [21-22]. Making informed decisions about maintenance energy usage and expling proformance optimization using the data is possible [23-25].

Data Lected:

\checkmark Solar PV Panels: P_{pv}

Power Output Data (P_{out}) : Each panel generates data on the amount of electrical power (P_{out}) it produces.

Environmental Data (E_{data}) : Solar panels are often equipped with sensors that collect environmental data (E_{data}), such as irradiance and temperature.

ii. DC/AC Inverter: $I_{dc/ac}$

Voltage and Current Data (V_{dc} , I_{dc} , V_{ac} , I_{ac}): The inverter captures data on the DC voltage (V_{dc}) and current (I_{dc}) from the PV panels and the AC voltage (V_{ac}) and current (I_{ac}) it outputs to the grid.

Efficiency Data (η_{inv}): It also records its operational efficiency (η_{inv}), measuring low we it converts DC to AC power.

iii. Generation Meter: M_{gen}

Energy Generated Data (E_{gen}): The generation meter logs the total vergy produced (E_{gen}), usually in kilowatt-hours (kWh).

iv. Consumer Unit: CU

Load Distribution Data (LD_{data}): The consumer unit provide data in load distribution (LD_{data}).

v. Electricity Meter: M_{elec}

Consumption Data (C_{data}): Records how nucle energy is consumed by the SG (C_{data}) and track energy exported back to the grid.

Net Usage Data (N_{usage}): Provides net usage ta (N_{usage}), which is the difference between energy produced and consumed.

vi. Monitoring and Control Systems: DC_{sys}

Performance Data (P_{data}) : These systems aggregate all the data (P_{data}) it provides a comprehensive overvew of the system's performance from various components.

Alerts and Fault entry is They also generate alerts and log fault data (A_{data}) .

vii. Qata communication: D_{com}

Transmission Data (TD_{data}): The system includes data transmission components that relay all collected data (TD_{data}) to a central monitoring point or off-site data centre for further analysis. *C. Letta Communication in Solar PV Power Plants*

In solar PV power plants, the usage of Supervisory Control and Data Acquisition (SCADA) systems is dynamic in the task of enabling remote management of several field devices, including sensors, smart meters, Remote Terminal Units (RTU), and Intelligent Electronic Devices (IED).

These systems comprise a network of components that coordinate to ensure operational efficiency and reliability.

The Components of a SCADA include:

- *Data Acquisition Units:* These are responsible for measuring and gathering key monitoring parameters like voltage, current, temperature, and irradiance.
- *RTU:* These units serve as the intermediary, collecting data from the acquisition processing it, and relaying it to the primary control system.
- *Communication Networks*: The system's backbone, facilitating data transfersion from the acquisition units to the control centre.
- *System Servers:* At the heart of the SCADA, many servers analy, an display the collected data for further action.

This includes:

- Front-End Servers: Tasked with aggregating data from the Fuld devices.
- *Historian Servers:* These servers act as data recositores, achiving data input for future reference and analysis.
- *Web Servers and Human-Machine Interfaces* (HMI): They provide a visual representation of the data and the status of the power plant, a pwing for real-time monitoring and control.

(i) Central Control Center:

The control center contains SCAPE, a which the application servers dissect the incoming information, making it users endly for perators to assess and act upon. This system performs data collection and viscon palyes functions to ensure that each PV power plant operates at its peak.

Illustrated in Fig. 2, the communication network of a PV showcases the interconnectivity between the local control centers, each dedicated to the management of a singular PV power plant. These centers are hacked through a vast area network, bridging the communication gap via routers, ensuring namless data flow and centralized control.

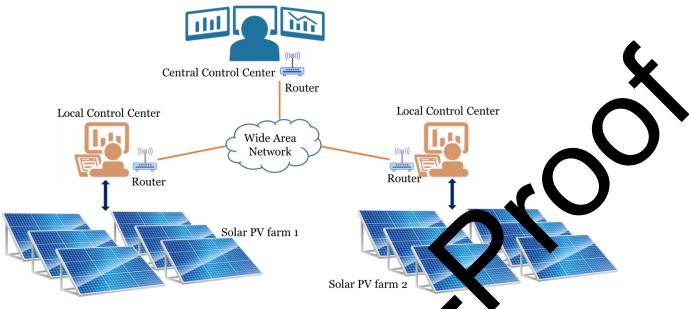


Figure 2. Communication network for a PV monit ring system.

(ii) Data and Power Integration in Communication Lauer

The data and power network is structured into the e-integral by yers, as shown in Fig. 3: the physical infrastructure of the PV, the data transmission backbone of the communication network, and the user-interfacing application layer.

a) The PV Power System Layer

The PV power system layer is the foundation of a large-scale solar plant. The physical layer is where sunlight is captured and opprove into usable electricity. This layer comprises all the essential equipment, including

PV Modules: The star cals that capture sunlight and convert it into electrical energy.

• Junction Boxe and Circuit Breakers: Safety devices that protect the system from electrical malfunctities.

• creection Devices: Equipment designed to shield the system from overloads and short circuits.

• *Fulnverters:* Devices that convert the DC generated by the PV modules into AC suitable for the popular grid.

Power Cables: Conductive wires that transport electricity throughout the plant.

• *Grid Connection Points:* Interfaces where the PV plant connects to the external power grid.

• *Transformers and Substations:* Apparatus that adjust voltage levels for efficient transmission and distribution.

Solar panels are connected in series to form a module string, which increases the voltage output. Multiple strings are then grouped and connected in a string combiner box. The string combiner boxes route the electricity to the Power Condition Unit (PCU), the system's main component. It converts the DC from the panels into AC. The specific type of transformer used depends on the overall plant design.

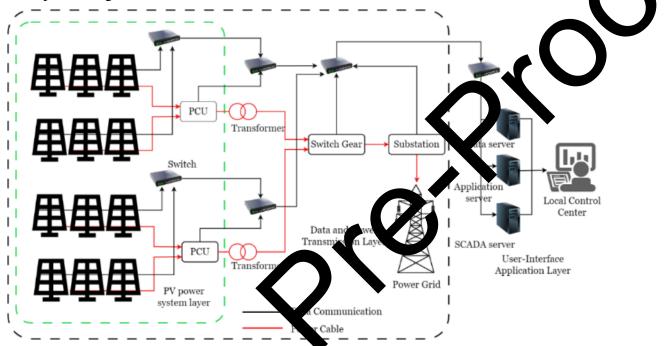


Figure 3. Layout of the integration of a large-scale PV power plant.

b) Communication Network Le

The communication network layer acts as the central nervous system of a large-scale PV power plant. It facilitates the critical two way flow of information between many subsystems and the local control centre.

This layer a porises everal key elements:

• *communication Devices:* A network of interconnected devices like cables, routers, and switch c ensures reliable data transmission throughout the plant.

Server Nodes and Measurement Devices: These intelligent devices gather data from the V processive system layer. They might monitor voltage, current, power output, ambient temperature, and speed.

• *Data Transmission:* The communication network layer transmits two main categories of data:

- *Monitoring Data from the PV subsystem:* This includes real-time information on the performance of various components, allowing for early detection of potential issues.
- *Meteorological Parameters:* Data on weather conditions, such as solar irradiance and wind speed, is crucial for optimizing energy production and maintenance schedules.

c) Application Layer

The application layer acts as its central brain. Here is how it transforms collected de intelligent control:

• Control Center Command: It receives data from the communication network layer, encompassing information on:

- PV Panel Performance
- Inverter Status
- Transformer Efficiency
- \circ Grid Connection Health

• *Data Analysis and Storage:* The control centrer oused server systems that collect, process, and store this data. And employ application to an ever use data to identify trends and potential issues and optimize performance.

• *Decision-Making and Control Actions.* The control center can make informed decisions and take appropriate actions based on the analyzed data.

This might involve:

- Adjusting inverter seeings to entiralize power output
- Activating main are potocols for faulty equipment
- Regulating power flow o meet grid requirements
- D. Attack Sectors and Cy. Threats in PV Power Systems

The breats on PV can be considered attack vectors that target PV's physical components and cyberinfrast octure. The Tab. 1 outlines the attack vectors and cyber threats specifically targeting PV power systems:

TABLE 1: ATTACKS/THREATS AND ITS IMPLICATIONS

	Engineering	Cruler of their; unauthorized system access
Interfaces	Phishing and Social	$\mathbf{\Omega}$
Software and Digital	Software Vulnerabilities	Exploitation leading to estem compromise; data
Internal Network	Insider Threats	Unintentional data contes; intentional sabotage
Smart Meters and Sensors	Protocol Exploits	Unauthorized access; data a anipulation
SCADA Smart Meters and Sensors	Firmware Tampering	Inaccurate data reporting, hergy heft
	Injection Attacks	System malfunction; unauthorized control
	Marware and Kansoni ware	for system control
Communication Networks	Malware and Ransomware	Manipulation of operational data; ransom daman
	Attacks	leading to system inefficiency
	Man-in-the-Middle (MitM)	Data tampering; incorrect operational command
	Attacks	operational downtime
	Denial of Service (DoS)	Disruption in monitoring and control capabilities;
	······································	access to sensitive information
Data Transmission	Data Interception and Theft	Compromise of data confidentiality; unauthorized

IV. PROPOSED ENDOGENOUS SECURITY IN PV

A. Assumptions

To validate the effectiveness of the propula endogenous CPS for PV data transmission and storage, we proceed under the following assumptions:

- *Resource Sufficiency:* The central controller (managing data flows and security protocols) and networked PV devices porcess ample processing capabilities and resources, ensuring minimal processing delays.
- *Bounded Propagation Dervy:* The propagation delay within the PV's communication network is limited, all using be central controller to monitor data packet movements accurately within acceptable time frames.
- Adequite Bardwidth: The communication links within the PV have sufficient bandwidth to minimize the packet loss due to network congestion, ensuring robust data transmission.
 - Secure Controller and Communication: The central controller and the control channel employed for data communication are secured using standard protocols such as Transport Layer Security (TLS), safeguarding against unauthorized access and data breaches.
- *Robust Encryption and Authentication:* The encryption and authentication mechanisms in place are considered secure against breaches (*e.g.*, encryption that cannot be easily broken and digital signatures that cannot be forged).

B. Threat Model

The endogenous security system is designed with the following threat models in mind for PV data transmission and storage systems:

- Focus on Core System Security: The primary concern is securing the core components of the PV's communication network. Threats related to peripheral or "edge" components, suc as individual PV modules or local inverters, fall outside the immediate scope of this research.
- Malicious Component Manipulation: There is a risk of malicious component within he network due to compromised hardware or software vulnerabilities the rest in the manipulation of data packets or flow rules.
- *Exclusion of Certain Network Attacks:* Attacks related to protocol, such as TCP/IP or OSPF, are considered beyond the scope of this model.
- C. Endogenous Security Model for PV Data Transmission

The endogenous security system for data transmission wintin PV uses enhanced routing strategies and a robust VFM to secure the data communication network. The model employs three routing approaches: (i) Data Integrity Forwarding DIF, which is for establishing accuracy and data consistency through data comparison communication by distributing data along different (ii) Load-Balanced Forwarding (LBF), handles the congestion by distributing data along different paths that were selected using predefined load capacity and (iii) Path Diversity Forwarding (PDF) is employed to select data transmission path from route pools to ensure randomness. The VFM employs the system to identify security attacks and anomalies by recognizing communication errors.

For processing, time management, and data path security and integrity validation, the structure includes the following sequents:

Enhanced with Computing (EPC): By examining the current network state and security hazards, Et a uses one of three algorithmic routing methods to determine the most efficient and secure routes for data transmission.

name: *Path Scheduling (DPS):* Implementing real-time variations in security and network unstances, this DPS unit objectives data transfers on demand.

Path Authentication and Feedback Verification (PAFV): This PAFV unit focuses on the data transfer route to identify malicious use behavior and information errors.

Data Path Validation Checking (DPVC): By verifying what is received to the implied structure,

this DPVC unit validates the packets of data integrity.

The operational flow of the processed model is presented below:

- 1. Upon initiating data transmission, the proposed system evaluates which routing strategy is most appropriate based on current network conditions and security protocols.
- 2. The controller then calculates the optimal paths for routing by employing the EPC respond to the network's needs adaptively.
- 3. Depending on the routing strategy selected:
- DIF routes duplicate data packets through different paths for cross-verification at the destination.
- LBF distributes data across available paths in alignment with the capacity to ensure a balanced load.
- *PDF* routes data packets via randomly selected paths o *cyustate* the transmission pattern and enhance security.
- 4. Data packets undergo a final verification process at their destination to confirm their integrity. Any detected irregularities are marked cancelalies.

The packet verification process is illustrated as follows:

- 1. *Initiation of Probe Packet:* The initial control system initiates the process by dispatching a probe packet.
- 2. *Verification Mark Generation:* A trice probe packet is transmitted, the first switch in its path generates two types of corks ased on the flow's characteristics and the packet's unique hash value:
 - 2.1. Ve if a tion Flow Fale Mark (e): This mark verifies flow against the predefined security uses a red on how information, flow entry point, and packet hash value.

Ver, "cation Data Content Mark (t): This mark validates the data content's integrity.

Intern diate Switch Processing: As the packet moves through subsequent switches in the

The verification information from the preceding switch (S_{i-1}) is encrypted with a key (K_i) and incorporated into the current switch's (S_i) Verification data. This layered approach ensures that each switch adds its unique verification mark to the packet.

- 3.2. Similarly, data verification information is cumulatively embedded and updated at each switch, with the current switch's data verification mark appended to the packet.
- 4. *Final Verification and Routing Decision:* Upon reaching the destination, the switch collects multipath information from various ports and forwards it to the control system. The control system then:
 - 4.1. A consistent comparing and ruling algorithm is used to evaluate the data from different paths. Key comparison metrics include the *Flow_ID*, *Datapath_ID*, *Buffer_D*, and he data message's hash value.
 - 4.2. The VFM scrutinizes the flow rules and data if discrepancies are detended. This facilitates precise identification and swift rectification of any anomalies in the paths.
 - 4.3. In cases where data paths need to be quickly changed, the system is designed to efficiently reroute specified data over alternate switch ports.

D. Enhanced Path Computing (EPC) and Dynamic Path Schedy (ng (DPS))

Constraints

The control system identifies the optimal data transmission paths across the PV by deploying discovery packets, which map the entire network in rout. Upon receipt of the initial data packet at a network node N_0 , this node signals the control stem, initiating the path optimization process. The objective of this path determination and optimization process is to select paths that minimize the number of intersecting nodes while satisfying specific transmission standards, EQU (1)

Minimize the intersection set $\sum_{(p_i, p_j) \in P, i < j} |I(p_i, p_j)|,$ (1)

where for every pair of pars p_{i} and p_{j} in the intersection set $I(p_{i}, p_{j})$ from source N_{src} to destination N_{dst} , each part c the flow paths subject to constraints:

Node Deversity equirement: For any selected group of paths within the possible flow sets, no single ode scauld be standard across all paths within the group, ensuring that compromised nodes cannot affect the entirety of the data transmission, EQU (2)

 $\forall p \in \square h \in p, "n_{\text{end}} " \notin p', \exists p' \in P \setminus p \tag{2}$

Bandwidth Limitation: The collective bandwidth of the selected multipath set must not exceed a predefined bandwidth threshold, ensuring adequate data flow without congestion, EQU (3) $\sum_{(p \in P, l \in p)} bw_p^l \leq B_{\text{total}}$ (3) *Link Integrity Assurance:* The operational status of each selected link is marked as 1, while unselected links are marked as 0. This does not account for potential link failures, EQU (4) $\forall l_{active} \in L, l_{active} = \{1 \text{ if } l \in p, 0 \text{ otherwise }\}$ (4)

Node Count Restriction: Within any selected path set, the count of nodes should not surpass the number of secure, operational nodes capable of proper data forwarding, EQU (5)

 $\forall g \in P, |N_g| \le |N_{\text{secure}}|, \text{ where } P \in \text{ Paths}$

Transmission Latency Bound: The cumulative transmission delay across any selected path mission exceed a maximum delay, ensuring timely data delivery, EQU (6)

$$\sum_{l \in p} t_l \leq T_{\max}$$
 , $\forall p \in P$

E. Data Path Validation and Selection

Upon data initiation at a network node N_0 , this node communicates the last flow details to the central control system. The control system logs the request stability and computes multiple data paths from N_0 to the destination node N_d , and selects paths matuminimize common points to enhance security and reliability. The transmission times an descense validation period for critical data flows where security outweighs latence, concluss. Data packets from N_0 are duplicated and dispatched across selected paths based on pada med rules. For example, replication forwarding might occur from N_0 to N_1 , N_2 , and N_3 by spectred actions. The decision on which paths to use considers the delay requirements, where residual packets directed consequently to ensure timely delivery to N_d .

Upon arrival at N_d , the databackets are aggregated, and a validation message is relayed to the control system. A corrensus mechanism validates the data, where packets sharing identical flow identifiers are considered. The network will consider these authenticated channels as secured routes for sus quere communications if the data is verified precisely. The system uses corrective methods are switching and route reconfiguration if differences evolve. The system will start a securit alerance all the required measures have been taken to deal with the errors. This warning will be used to clean up and repair affected routes.

thentication and VFM

F. Pa.

The following methods are used to develop this system to identify and mitigate assaults.

- 1. Key Distribution and Flow Identification:
- Utilizing asymmetric cryptography, the controller maintains a key pair (K_{public} , K_{private}) and distributes individual session keys (K_{session}) to network nodes (N_i).

A unique flow identifier (FID) is generated for each data flow based on its characteristics:
FID = HashFn(PortEntry || FlowHeader)

Here, *HashFn* denotes a secure hash function, PortEntry is the input path for the data flow, and Flow Header includes dangerous header details like source and destination identifiers.

- 2. Probe Packet Dispatch and Response:
- When the starting node N_0 meets data *d* missing a match, it computes Hash(d), for robust this along with *d* to the controller. The controller, in response, techniques a payloid encompassing *FID*, a flow header, a protection timestamp *TS*, and a second starting $Sig_{K_{private}}$ (*FID* || *TS*) to ensure authenticity and temporal integrit
- The first node N_0 then propagates d along with its Hash (d), FID, N and $S_{ig_{K_{prinate}}}$ to the next node in the sequence, N_1 .
- 3. Validation and Encryption at Intermediate Nodes:
- At each node N_i , based on a predefined Validation (ag (VF), the data may undergo authentication to verify the integrity of the calculated signature and the associated time stamp, ensuring that the VF has not been altered.
- An encryption function at the node N_1 appends an authentication tag Tag_{N_1} to d, computed as: $\operatorname{Tag}_{N_1} = \operatorname{MAC} C_{K_{\text{session}}}$ (PortEntry $(N_1) \parallel$ AunPayload) where AuthPayload = FID $\parallel TS$, and MAC is the message authenticated code generated using the session key.
- 4. Data Forwarding and Corprehensive Verification:
- Successive nodes the tensine the necessity for authentication via VF, continuously appending verification tage to ensure end-to-end integrity.
- Upon erriver as the dominal node N_n , a consolidated packet containing d, Tag_{N_n} , and Hase (d) is dispatched to the controller for final validation, employing a comprehensive check against up original flow details.

Contr Ver's Final Authentication and Feedback:

- the controller executes a thorough verification for packets sharing an *FID*, leveraging the sion keys to authenticate the flow's transmission path and data content.
- This step facilitates precisely identifying anomalous or compromised nodes, enabling swift corrective actions to reestablish secure data pathways.

The following algorithm presents the steps in detail about the proposed security architecture.

Output: Secure and optimized data transmission paths Initialize keyPair (K_{public} , K_{private}) for the controller. Distribute K_{session} to all network nodes N_i . Initialize FlowTable as an empty dictionary. Enhanced Path Computing (EPC) 1. For Each data request *DR* in Data requests, **Do** 1.1. $FID = HashFn(DR.PortEntry \parallel DR.FlowHeader)$ 1.2. Paths = DiscoverPaths(DR.Source, DR.Destination)1.3. OptimalPaths = ComputeOptimalPaths (Paths) 1.4. FlowTable [FID] = OptimalPathsDynamic Path Scheduling (DPS) 2. For Each FID in FlowTable.keys() Do 2.1. SelectedPath = SelectPathBa (FlowTable [FID]) dOnS ateg 2.2. If strategy == DIF then 2.2.1. PerformDIF(SelectedPath) 2.3. Else, if strategy = LP 2.3.1. PerformLBF 2.4. Else nPD, (Se. stedPath) 2.4.1. Perfo Path Authentication and Feed ack Verification (PAFV) ath in SelectedPath, Do 3. For Ea $ober cket = GenerateProbePacket (K_{public}, K_{private}, FID)$ Send reperence Packet through the path. 3. For Each node N_i in path **Do** 5.1. Tag N_i = AuthenticateNode (N_i , ProbePacket, $K_{session}$) 3.3.2. Append Tag_{N_i} to ProbePacket. 3.4. Feedback = CollectFeedback (ProbePacket) 3.5. If VerifyFeedback(Feedback) is False then

Algorithm 1: Endogenous Security for PV Data Transmission

Input: Network topology, Data requests

1

2

3

- 3.5.1. Alert and recompute OptimalPaths excluding compromised path.
- 4. For Each path in SelectedPath, Do
 - 4.1. DataTransmission(path)
- (i) Function PerformDIF(path)
- Copy and forward packets through different paths
- Cross-verify at destination
- Forward only if data is consistent
- (ii) Function PerformLBF(path)
- Distribute data across paths based on capacity
- Ensure balanced load
- (iii) Function PerformPDF(path)
- Select paths randomly
- Increase unpredictability for attackers
- (iv) Function DiscoverPaths(source, destinction)
- Use LLDP or a similar protocol to discover all ossible paths
- Return list of paths
- (v) Function ComputeOptimalPaths(paths)
- Apply heuristic algorithms to f ad optimal paths minimizing node intersection
- Return optimal paths
- (vi) Function SelectPathBase OnStrategy(paths)
 - Determine strengy band on current network conditions and security needs
 - Return the electric ding of for data transmission
- G. Key Mana, ment a Distribution

Let $Roc(K_{pub}, K_{priv})$ represent the RSA algorithm generating public key (PuK) and private key (Phi) panel or asymmetric encryption, where K_{pub} denotes the PuK and K_{priv} signifies the conceptoning PrK. Similarly, $AES(K_{sym})$ symbolizes the generation of symmetric keys (SyK) is using the AES-256 standard, with K_{sym} indicating the symmetric key.

Key distribution is facilitated through a secure channel, denoted as SC, which employs Transport Layer Security (TLS) protocols to ensure the confidential transfer of K_{sym} to network components. This process is as $SC_{TLS}(D_i, K_{sym})$, where D_i is the *i*th device in the network receiving its SyK, K_{sym} .

The key management is integrated into CSP as follows:

DIF: For each data packet P_{data} encryption is applied using the symmetric ker $Enc_{AES}(P_{data}, K_{sym})$ before transmission, this encrypted data ensures that integrity checks during DIF are performed on secure content, thereby preserving data confidentiality and interviry ross the transmission paths.

LBF and PDF: The selection of paths for LBF and PDF is predication in availability of secure communication channels. Let P_{opt} represent the optimal path elected brough the heuristic evaluation of path security and network conditions, formulated as Selecteath (D_i, K_{sym} , NetState) $\rightarrow P_{opt}$, where NetState encapsulates the current network state, including load and security posture.

VFM: The integrity of feedback and probe pacters, k_{edback} , is secured through digital signatures Sig_{RSA} ($P_{feedback}$, K_{priv}), ensuring the a them sity and non-repudiation of the feedback sent to the control system for anomaly detection and system adjustments.

The Key Management System (KMS) automates the processes of key generation, distribution, rotation, and revocation within the model as $KMS(K_{pub}, K_{priv}, K_{sym}, SC_{TLS})$. This system ensures that cryptographic keys are dynamical, managed in response to network events, security incidents, or predefined schedules, enterting the resilience of the communication infrastructure. V. IMPLEMENTATION AND CALCUTION

A. Implementation Degils

Hardware Lawin open. The hardware setup consists of commercial-grade solar panels rated at 300 W geak power connected to a central inverter with a capacity of 10 kW. The network infrastructur is built using Cisco Catalyst 2960-X Series Switches. Each component within the PV is neurfaced with Raspberry Pi-4 Model B devices for real-time data processing and encention

Environment: The control system software is developed on the Node-RED platform, and the security model is encoded using Python 3.8 and Scapy libraries for packet manipulation and PyCryptoDome for cryptographic functions. For simulation, the NS-3 network simulator is employed. Table 1. presenting the configuration of the system setup:

TABLE II. SYSTEM CONFIGURATION

Component	Specification/Tool	Description/Function
PV Modules	300W peak power	High-efficiency solar panels capturing SE.
Inverter	10kW capacity	Converts DC to AC power, equipped with network
		interfaces for secure data communication.
Data Loggers	Raspberry Pi 4 Model B	Collects and records system performance data, equip
		with secure transmission capabilities
Network Devices	Cisco Catalyst 2960-X	Facilitates encrypted data communication whin the P
	Switches	system.
Control System	Node-RED platform	Manages monitoring, data ow, ed im, ementation of
		security algorithms.
Security Software	Python 3.8, Scapy,	Executes the endogenous durity model, including
	PyCryptoDome	cryptographic functions and cket manipulation.
Simulation Tools	NS-3	Models the PV and nework infrastructure for testing
		under simulated conditions.

B. Performance Metrics

The effectiveness of the endogenous 6.5 Equalitatively assessed using the following metrics, each represented with corresponding formulas to ensure precise evaluation, EQU (7) to EQU (10).

Data Transmission Integrity (I) :

$$I = \frac{N_{\text{correct}}}{N_{\text{total}}} \times 100\%$$

where N_{correct} is the number of data packets received accurately at the destination, and N_{total} is the total number of data packets ent. System Resilience to Attacks (A.):

(8)

(7)

here, T_{attack} represents NT under attack conditions and T_{normal} signifies system NT under normal operatio

(9)

where D_{total} is the total data transmitted over the network in a specific timeframe, and t is the time taken.

Latency (L) :

R = 1

$L = t_{\text{destination}} - t_{\text{source}}$

with $t_{destination}$ being the time a packet is received at its destination and t_{source} the time it was sent from its source. This measures the delay introduced by security protocols.

The above metrics were compared against the attacks such as (i) Data Flow Manipulation, (ii) DoS Attacks, (iii) Spoofing Attacks, and (iv) Path Compromise. The proposed models' performance was compared against the works, and the results analysis for the above metrics a discussed below:

The transmission integrity analysis of the compared models was done using different necket sizes and attacks. The Fig. 4 shows the transmission integrity of the compared model for different types of attacks. The proposed model shows better transmission integrity scales ranging from 97% to 99% compared to other models. The compared models show limited performance, particularly for DoS and Spoofing attacks, for which the proposed model displayed better performance.

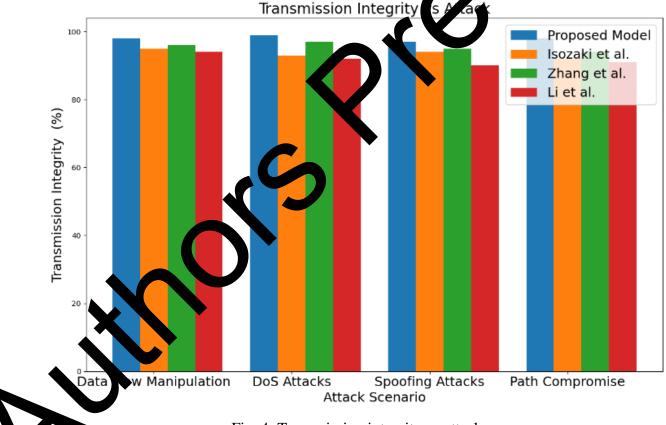
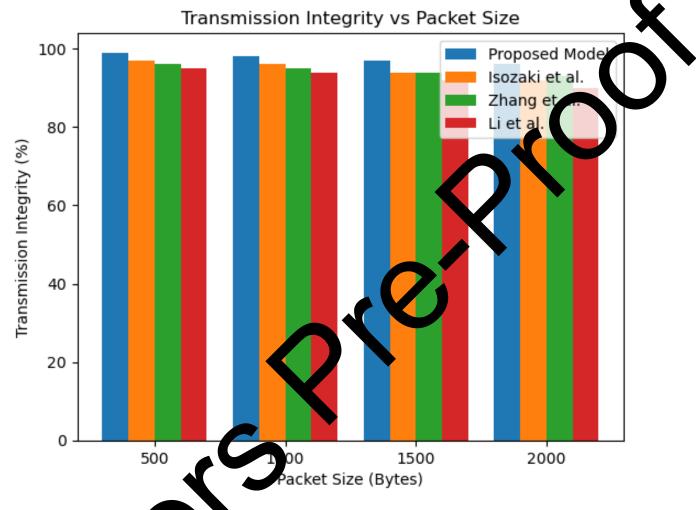


Fig. 4: Transmission integrity vs attack

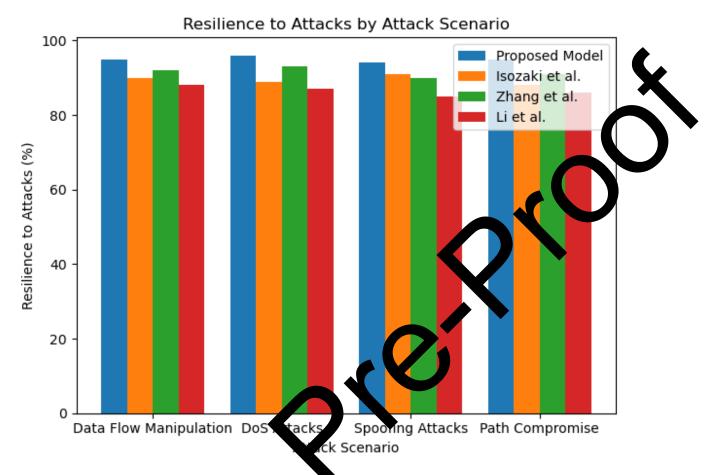
The Fig. 5 shows the comparison of transmission integrity against different packet sizes. The proposed model performs better with 96% and 99% integrity rates for all packet sizes tested

ranging from 500 to 2000 *bytes*. The performance decreased as the packet size increased, which was visible across all models, but even then, the proposed model showed better performance.



5. Cansmission integrity vs Packet size

The analysis of "Re lience o Attacks" (*R*) across different security models is shown in Fig. 6. For Data F ulation and Path Compromise attacks, the proposed model maintained a Man 25%, which is close to normal conditions. For DoS Attacks, the proposed model resilien ate o resilience rate of 96%, which is a better rate considering it is a high-intensity threat. shows high g Attacks, the proposed model had shown a resilience rate of 94%. Among the Spool comp hodels, Isozaki et al. showed resilience rates ranging from 88% to 91%, with the lowest resinchce for Path Compromise attacks and the highest for Spoofing Attacks. Zhang et al.'s model showed resilience rates between 90% and 93%, which was balanced compared to the other models except the proposed model.



Fig_6. Resilience vs Attacks

roue (T) across various security models is shown in Fig. The evaluation of "Network T proposed nodel recorded an NT of 10 Gbps, which is better than 7. Under normal operation, other models with Zhap ગ.'ડ odel is the next model reaching 9.8 Gbps. For DoS Attacks, the NT for the proposed odel decreases to 9.2 Gbps, which is followed by Zhang et al.'s model that of Spoofing attacks, the proposed model achieved NT of 9.6 Gbps, achieved 4 G th ppromise situations, the proposed model demonstrates an NT of 9.7 Gbps. Out of and for Rath s, the one next to the proposed model was Zhang et al., and Li et al. was the leastall the mos perform g model.

the ant lysis of "Latency" (L) is presented in Fig. 8. In normal operations, the proposed model bibits the lowest latency of 20 *ms*, Zhang et al. have shown a performance of 22 *ms*, and Isozaki et al. and Li et al. had shown higher latency of 25 and 28 *ms* respectively. For data flow manipulation, the proposed model shows 25, followed by Zhang et al. at 27 ms and Isozaki et al. at 30 ms, which comes at the last. A similar trend is understood across all attacks; the model had shown 30 ms latency for DoS attacks. For Spoofing Attacks, the proposed model shows lower latency at 23 *ms;* for Path Compromise attacks, the proposed model exhibits a latency of 22 *ms*. Among all models, Zhang et al. come next to the proposed model, and Li et al. scored the lowest performance across all attacks.

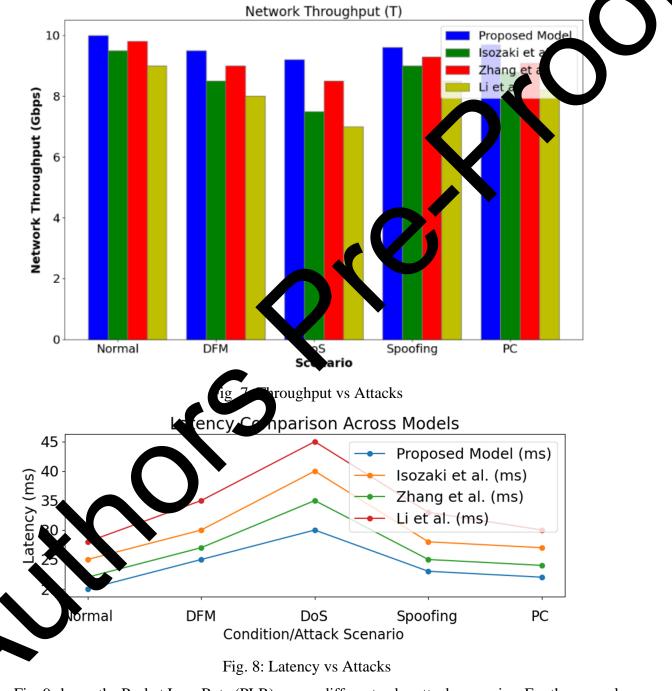
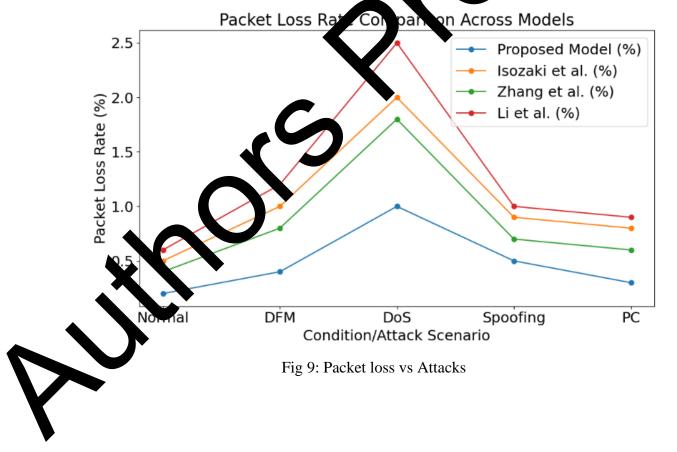
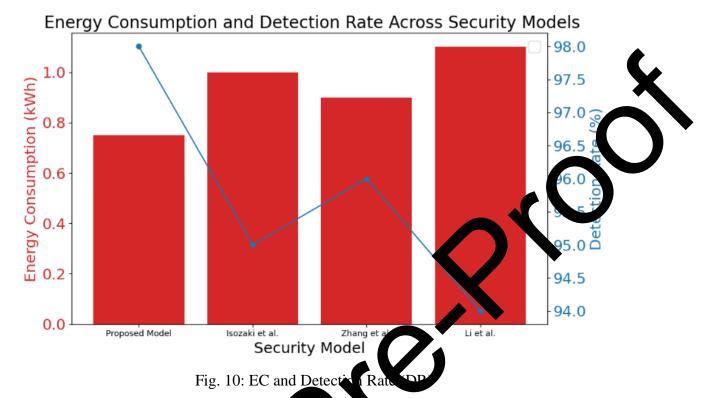


Fig. 9 shows the Packet Loss Rate (PLR) across different cyberattack scenarios. For the normal operation, the proposed model exhibits a packet loss rate of 0.2%, the lowest among the models,

indicating high data transmission reliability. Isozaki et al. record 0.5%, Zhang et al. 0.4%, and Li et al. 0.6%. The proposed model maintains the PLR at 0.4% for the data flow manipulation. Isozaki et al. have a rate of 1.0%, Zhang et al. 0.8%, and Li et al. the highest at 1.2%. For the DoS attacks, the rates increase due to the attack's nature, with the proposed model at 1.0%, demonstrating resilience. Isozaki et al.'s PLR is 2.0%, Zhang et al.'s 1.8%, and Li et al.'s 2.5%. As for the spooling attacks, the proposed model shows a 0.5% PLR, compared to Isozaki et al. at 0.9%, Zheng et al. 0.7%, and Li et al. at 1.0%. The proposed model records a 0.3% PLR for the Path Compromise, indicating effective rerouting strategies. Isozaki et al. have 0.8%, Zhang et al. 0.4%, and Li et al. 0.9%.

The EC and DR accuracy are analyzed in Fig. 10. The proposed model domonstrates the lowest EC at 0.75 kWh. Zhang et al.'s model follows at 0.90 kWh, Isozaki et al. model consumes 1.00 kWh, and Li et al.'s model has the highest at 1.10 kWh. The proposed model achieves a 98% DR for the analysis, the highest among the compared models. Inany et al.'s model has a DR of 96%, Isozaki et al.'s model records a 95% DR, and Li et al.'s model has the lowest at 94%.





VI. CONCLUSION AND FUTURE WORK

Due to weaknesses in digitalized V overations, the move toward renewable energy sources (RES) has made protecting critical systems from cyberattacks more important. Using routing methods like DIF, LBF, and PDF, this research has developed an autonomous security system that protects data privacy which has being sent and stored. A cryptographic key system improves network communication placey. This study recommends models for data integrity, system reliability, latency, and hetwork traffic. The proposed model does better in all of these ranges than competing models

Funce esea h should explore machine learning (ML) for predictive threat detection and expand by morel to include numerous RES.

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